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Basic concepts

We are given with

• a table is an array $a = [a_i : i \in I]$ satisfying:

$$\sum_{i \in I} m_{ij} y_i = b_j \quad \text{for all } j \in J$$
$$lb_i \le y_i \le ub_i \quad \text{for all } i \in I.$$

• A subset of sensitive cells $P \subset I$ with $[LPL_p, UPL_p]$ for each $p \in P$.

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- A subset of sensitive cells $P \subset I$ with $[LPL_p, UPL_p]$ for each $p \in P$.
- Find "something" such that a data user will see the original table and other tables as possible tables such that:
 - LPL_p and UPL_p should be possible values in a table.
 - The loss of information due to the existence of other tables is minimized.

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- Alternative: Danderkar, R.A. and Cox, L.H. (2002). "Synthetic Tabular Data - An Alternative to Complementary Cell Suppression", manuscript. Energy Information Administration, U.S. Department of Energy.

Cell Suppression: mathematical model



Subject to:

$$\begin{array}{ll} \underline{y}_p \leq lpl_p &, \quad \overline{y}_p \geq upl_p &, \quad \overline{y}_p - \underline{y}_p \geq spl_p &\quad \forall p \in P \\ & x_i \in \{0, 1\} &\quad \forall i \in I \setminus P \end{array}$$

where

$$\underline{y}_{p} := \min y_{p} \quad \text{and} \quad \overline{y}_{p} := \max y_{p}$$

$$\sum_{i \in I} m_{ij} y_{i} = b_{j} \quad \text{for all } j \in J$$

$$a_{i} - (a_{i} - lb_{i}) x_{i} \leq y_{i} \leq a_{i} + (ub_{i} - a_{i}) x_{i} \quad \text{for all } i \in I$$

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$$\begin{array}{ll} \min \sum_{i \in I} w_i | y_i - a_i | \\ \text{subject to:} & \sum_{i \in I} m_{ij} y_i = b_j & \text{for all } j \in J \\ lb_i \leq y_i \leq ub_i & \text{for all } i \in I \\ y_p \leq lpl_i & \text{or } y_p \geq upl_i & \text{for all } p \in P \end{array}$$

which one has smaller loss of information?

				_				
34566	—	—	92525		_	_	_	92525
53453	66345	43563	163361		_	_	_	163361
145343	_	_	243131		_	_	_	243131
233362	113315	152340	499017		233362	113315	152340	499017

which one has smaller loss of information?

34566	3425	54534	92525	66477	5730	20318	92525
53453	66345	43563	163361	89552	53242	20567	163361
145343	43545	54243	243131	77333	54343	111455	243131
233362	113315	152340	499017	233362	113315	152340	499017

Motivation:

- Utility: How useful is a perturbed data from the user point of view?
- Protection: How can we ensure both lower AND upper protection levels on each sensitive cell?

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Motivation:

- Utility: How useful is a perturbed data from the user point of view?
- Protection: How can we ensure both lower AND upper protection levels on each sensitive cell?
- Complexity: CTA, like CS, requires solving an Integer Linear Programming model, thus their optimization problems are both NP-hard.
- Instead, ECTA consists in solving a model without binary variables, with only 1 variable y_i for cell $i \in I$, and with only 1 additional variable β that will be also published.

- **Select** k sensitive cells in a random way.
- Fix each of these cells to a random value ξ_p in $[lpl_p, ulp_p]$.
- Fix the other sensitive cells to their original value $\xi_p := a_p$.
- **Solve the LP model:** $\min \beta$

subject to
$$\sum_{i \in I} m_{ij} y_i = b_j$$
 for all $j \in J$
 $lb_i \leq y_i \leq ub_i$ for all $i \in I$
 $y_p = \xi_p$ for all $p \in P$
 $\left(1 - \frac{\beta}{2}\right) a_i \leq y_i \leq \left(1 + \frac{\beta}{2}\right) a_i$ for all $i \in I \setminus P$

If the problem is infeasible or the solution is unprotected then restart with a larger k.

ECTA: computational results

- we have implemented CTA and ECTA using a free-and-open-source optimizer: SCIP.
- we considered 2 instances from the CSPLIB where this solver needs more than 1 hour: Hier16, Ninenew

Ninonowcor

Hior16 con

ECTA found protected solutions in a few minutes

		niei io.csp	Ninenew.csp	
	I	3564	6546	
	P	224	858	
	J	5484	7340	
	Number of ECTA models	30	10	
Details: $(k+=5)$	Total ECTA time	197 seconds	459 secs	
	Time for finding (eta,y)	51 secs	72 secs	
	Time for auditing (eta,y)	143 secs	382 secs	
	Protected solutions	1 sol	1 sol	
	Non-protected solutions	29 sols	9 sols	
	Infeasible solutions	0 sols		stment – p. 10/10